

Asymmetric Key Encryption

- Important to know who should know which key(s)
- In general:
 - Sender encrypts with recipient's public key
 - Recipient decrypts with its private key

Matrix of Keys

Key details	<i>A</i> should know	<i>B</i> should know
A's private key	Yes	No
A's public key	Yes	Yes
B's private key	No	Yes
B's public key	Yes	Yes

Fig 4.1

Asymmetric Key Cryptography

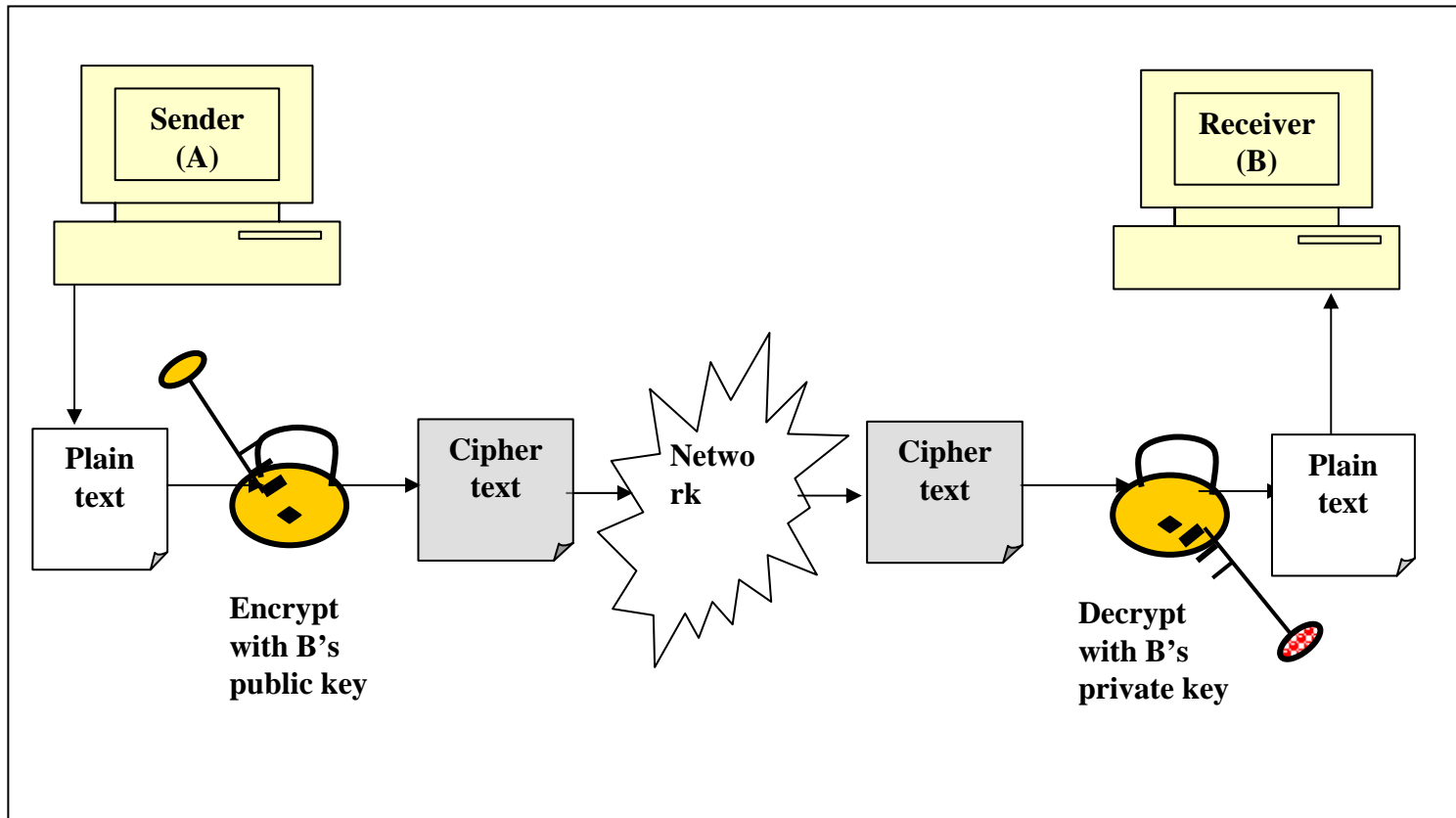


Fig 4.2

Asymmetric Key Example

- Consider a bank and its customers
- Customers encrypt their messages with bank's public key
- Bank decrypts messages with its private key

Asymmetric Key Cryptography Example

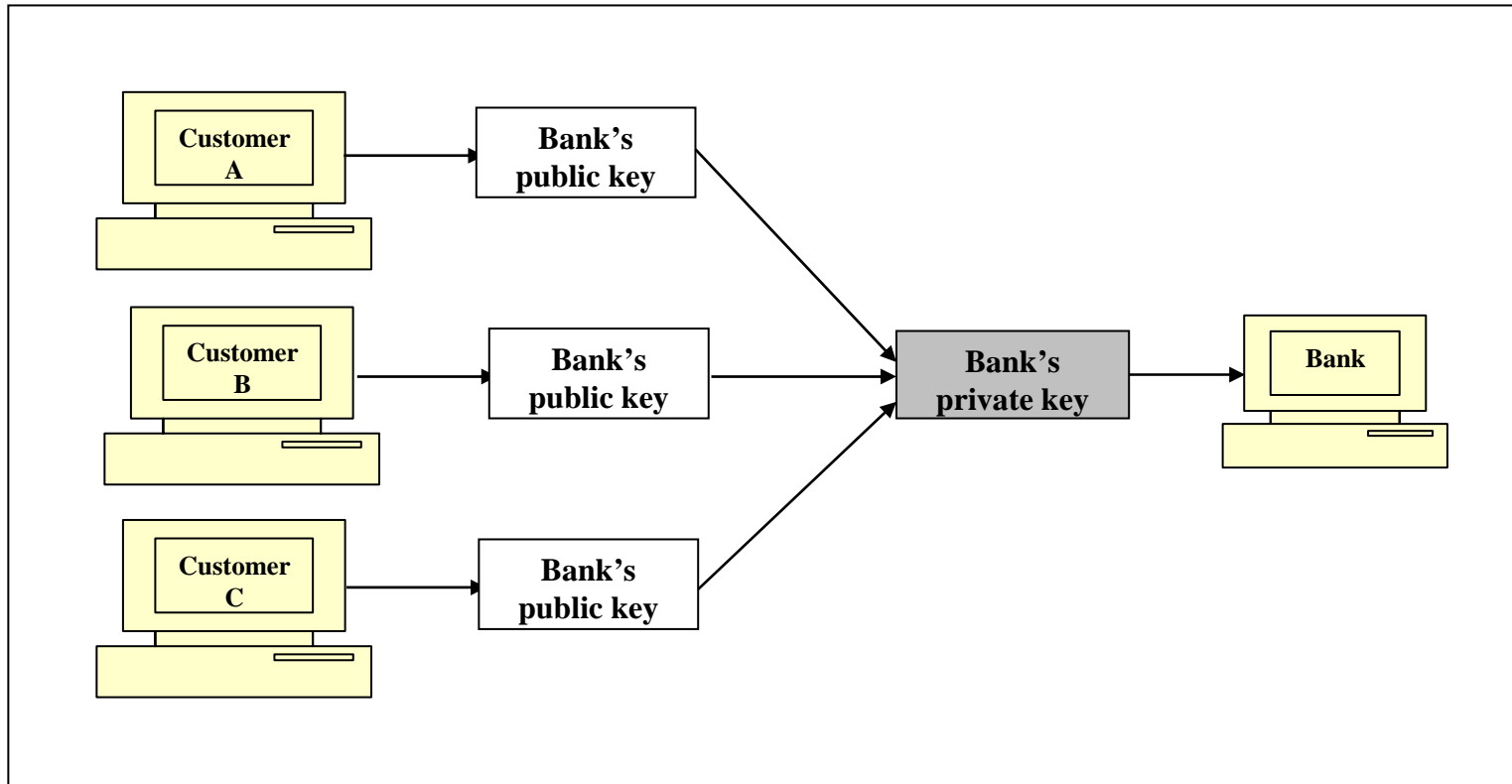


Fig 4.3

RSA Algorithm

- World's most popular Asymmetric Key Encryption algorithm
- Based on the theory of Prime Numbers
- Example Follows

RSA Algorithm

- 1. Choose two large prime numbers P and Q.**
- 2. Calculate $N = P \times Q$.**
- 3. Select the public key (i.e. the encryption key) E such that it is not a factor of $(P - 1)$ and $(Q - 1)$.**
- 4. Select the private key (i.e. the decryption key) D such that the following equation is true:
 $(D \times E) \bmod (P - 1) \times (Q - 1) = 1$**
- 5. For encryption, calculate the cipher text CT from the plain text PT as follows:
 $CT = PT^E \bmod N$**
- 6. Send CT as the cipher text to the receiver.**
- 7. For decryption, calculate the plain text PT from the cipher text CT as follows:
 $PT = CT^D \bmod N$**

Fig 4.4

Example of RSA Algorithm

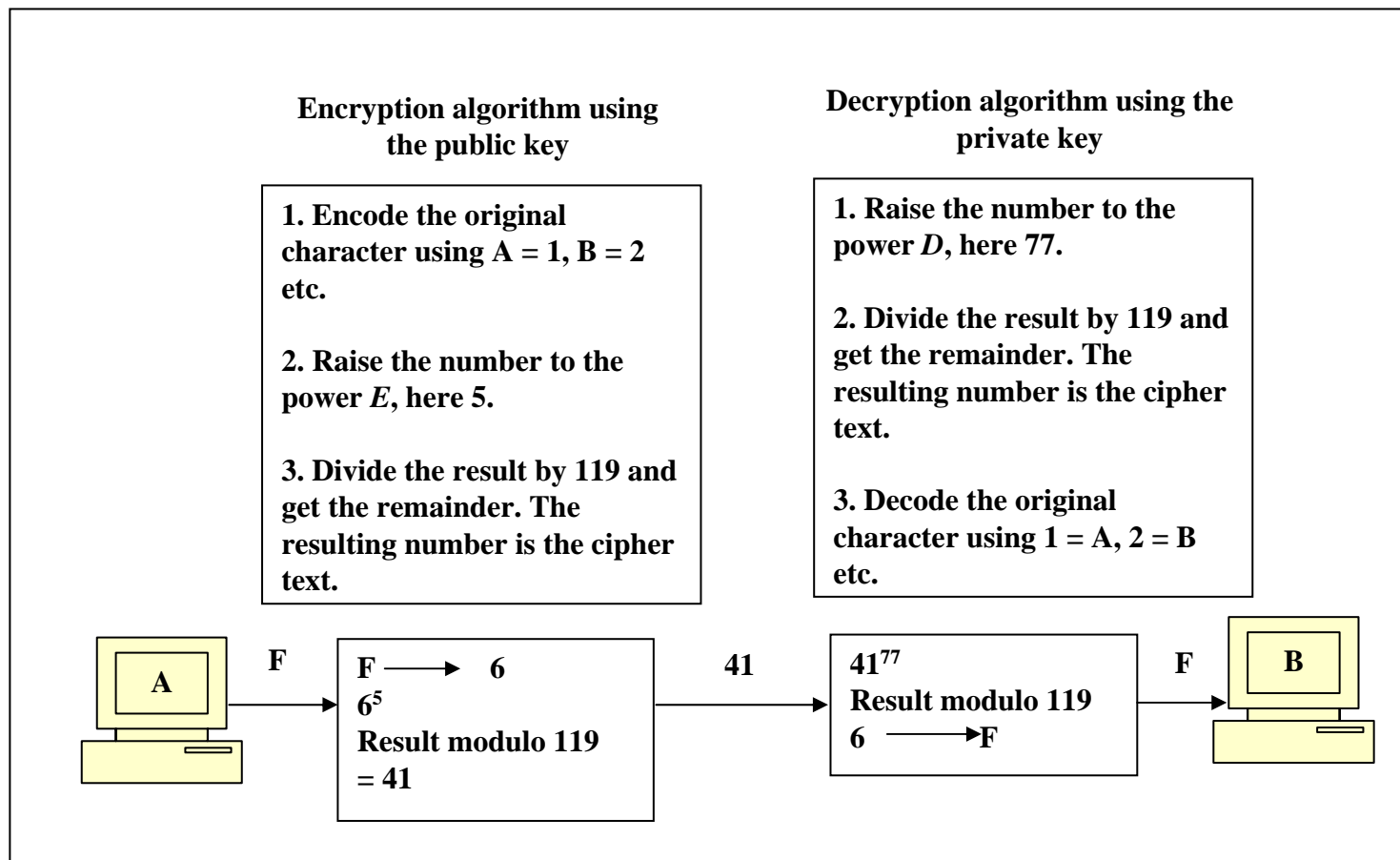


Fig 4.6

Symmetric v/s Asymmetric

Characteristic	Symmetric Key Cryptography	Asymmetric Key Cryptography
Key used for encryption / decryption	Same key is used for encryption and decryption	One key used for encryption and another, different key is used for decryption
Speed of encryption / decryption	Very fast	Slower
Size of resulting encrypted text	Usually same as or less than the original clear text size	More than the original clear text size
Key agreement / exchange	A big problem	No problem at all
Number of keys required as compared to the number of participants in the message exchange	Equals about the square of the number of participants, so scalability is an issue	Same as the number of participants, so scales up quite well
Usage	Mainly for encryption and decryption (confidentiality), cannot be used for digital signatures (integrity and non-repudiation checks)	Can be used for encryption and decryption (confidentiality) as well as for digital signatures (integrity and non-repudiation checks)

Fig 4.7

Digital Signature Concept

- Sender encrypts message or its fingerprint with its private key
- Guarantees that only the sender could have created this message
- Basis for Non-repudiation

Basis for Digital Signatures

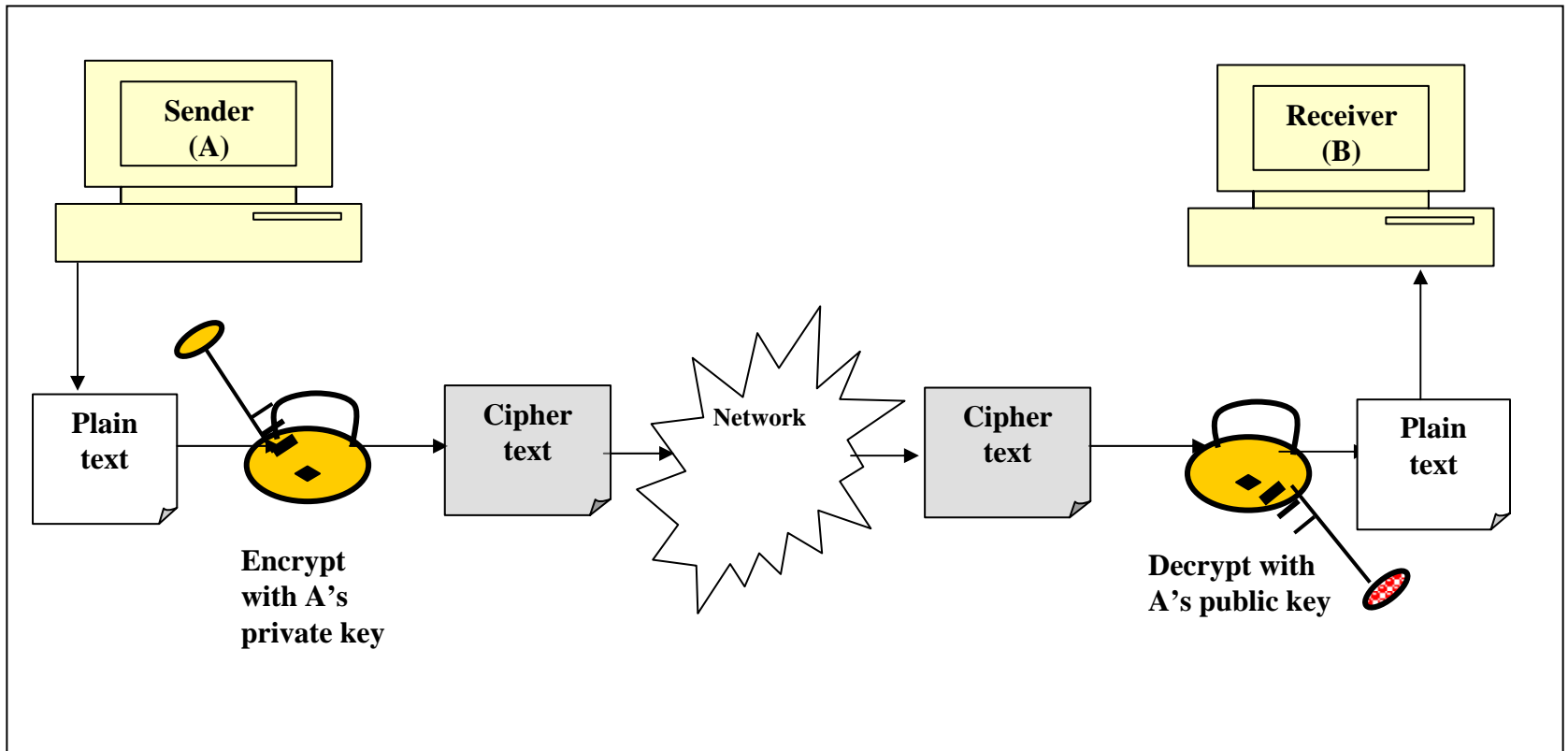


Fig 4.16

Message Digest Concept

- Also called as *Hash*
- Unique representation of a message
- Similar to finger print of a human

Message Digest Idea

•Original number is 7391743

Operation	Result
Multiply 7 by 3	21
Discard first digit	1
Multiply 1 by 9	9
Multiply 9 by 1	9
Multiply 9 by 7	63
Discard first digit	3
Multiply 3 by 4	12
Discard first digit	2
Multiply 2 by 3	6

•Message digest is 6

Fig 4.18

Message Digest Concept

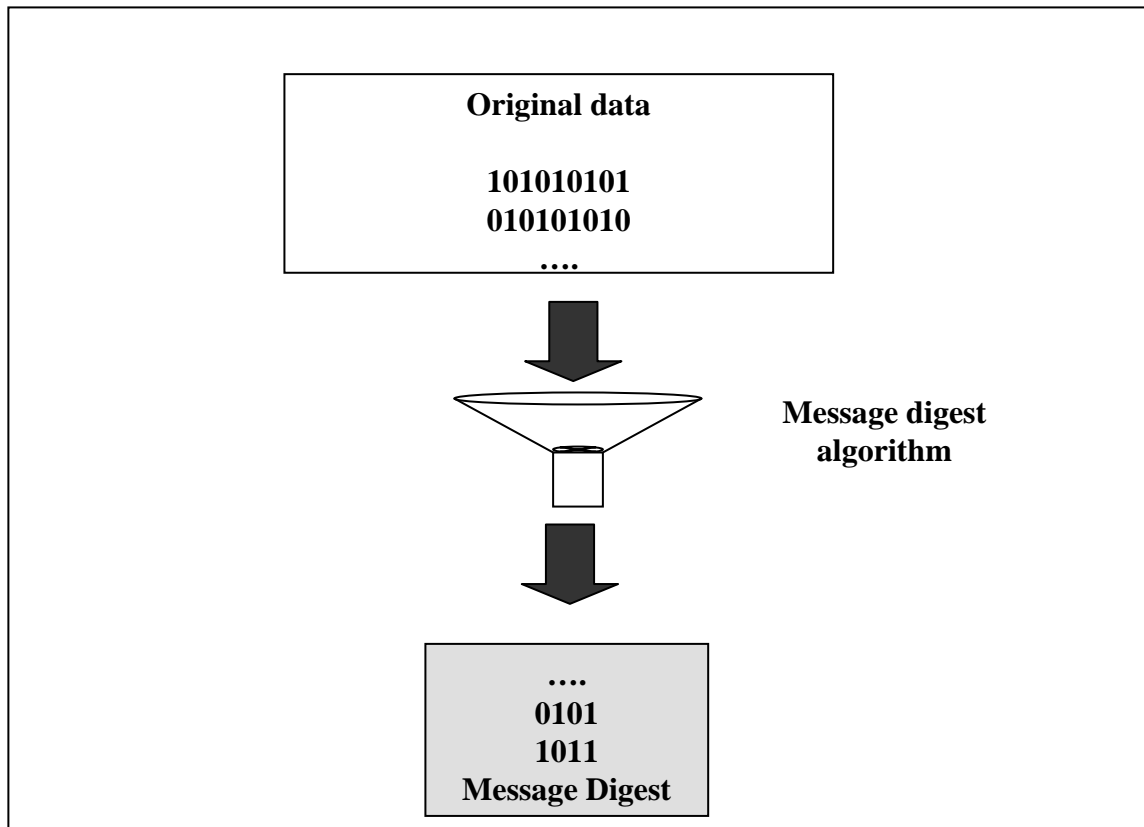


Fig 4.19

Message Digest Demands - 1

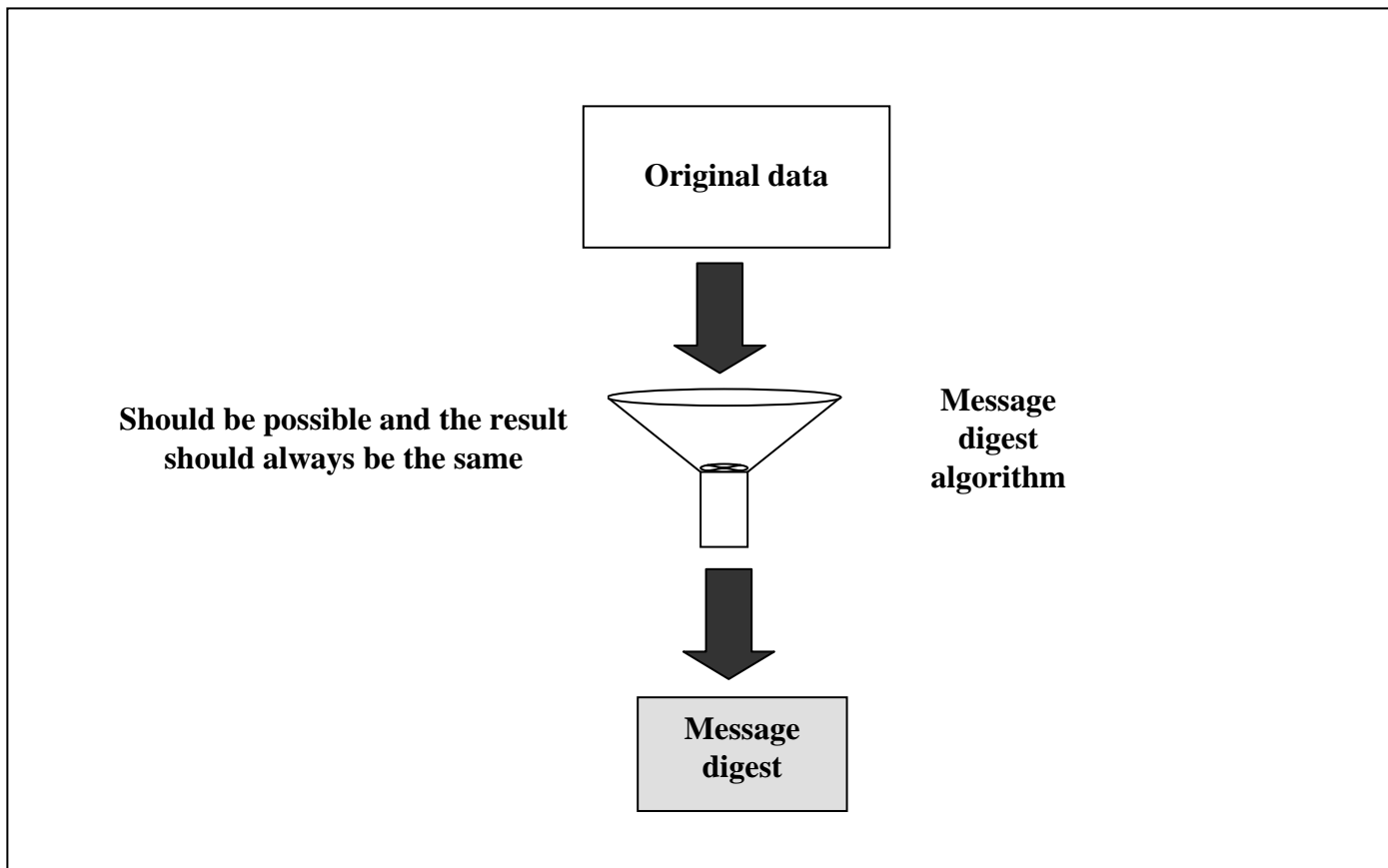


Fig 4.20

Message Digest Demands - 2

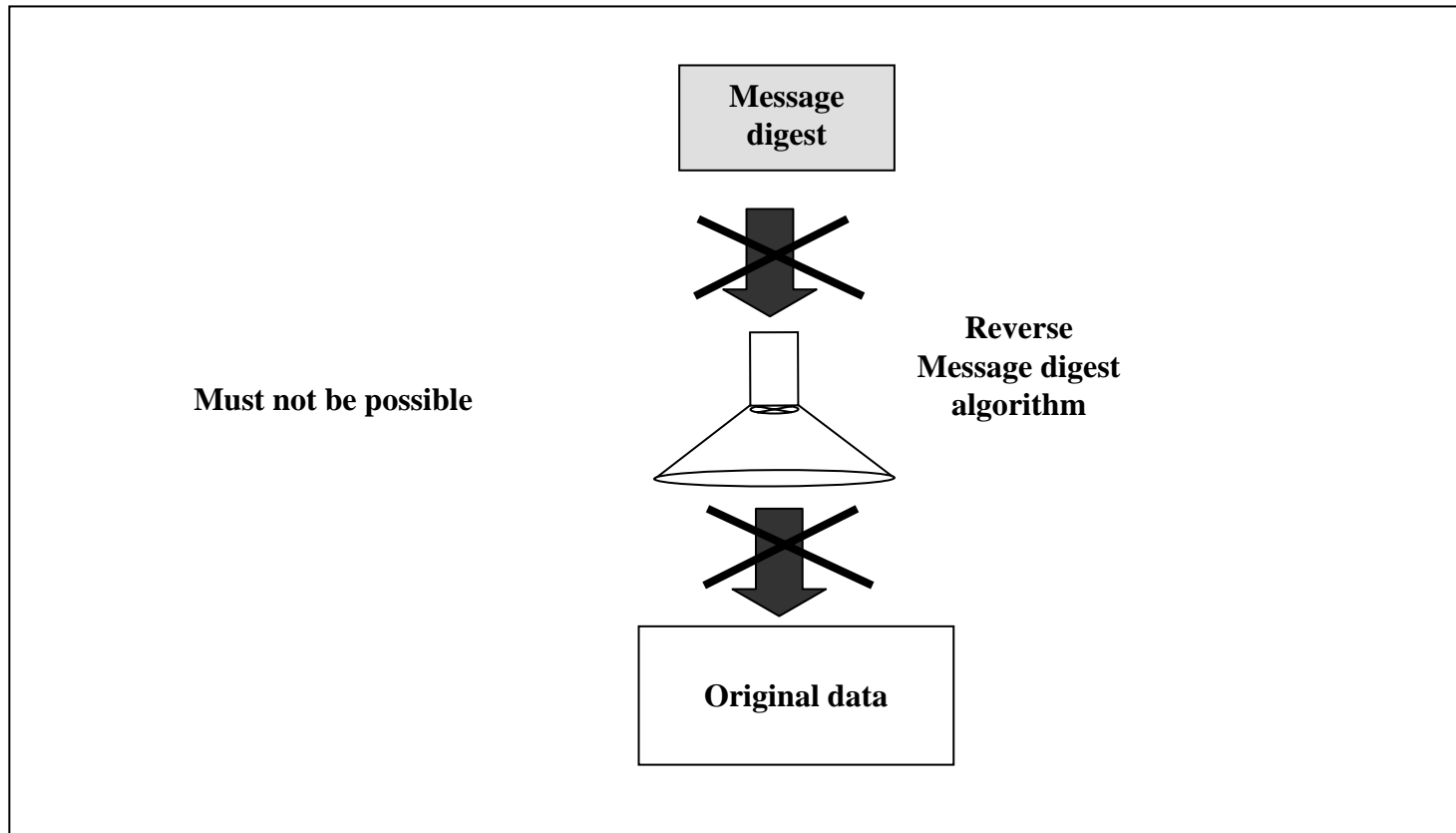


Fig 4.21

Message Digest Demands - 3

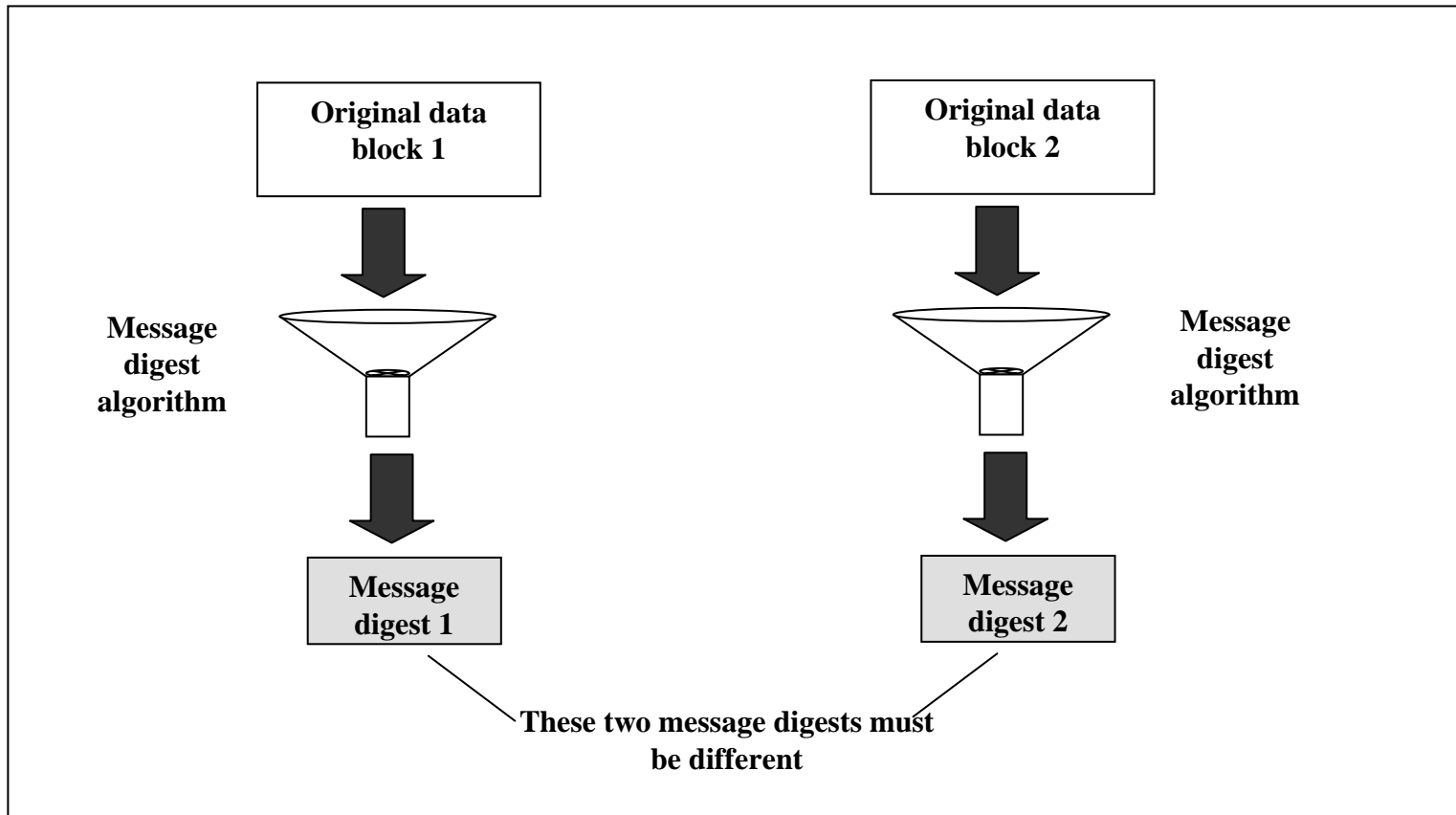


Fig 4.22

Message Digest Differences

- Even if the original messages differ minutely, message digests differ dramatically
- Basis for the guarantee of uniqueness

Message Digest Example

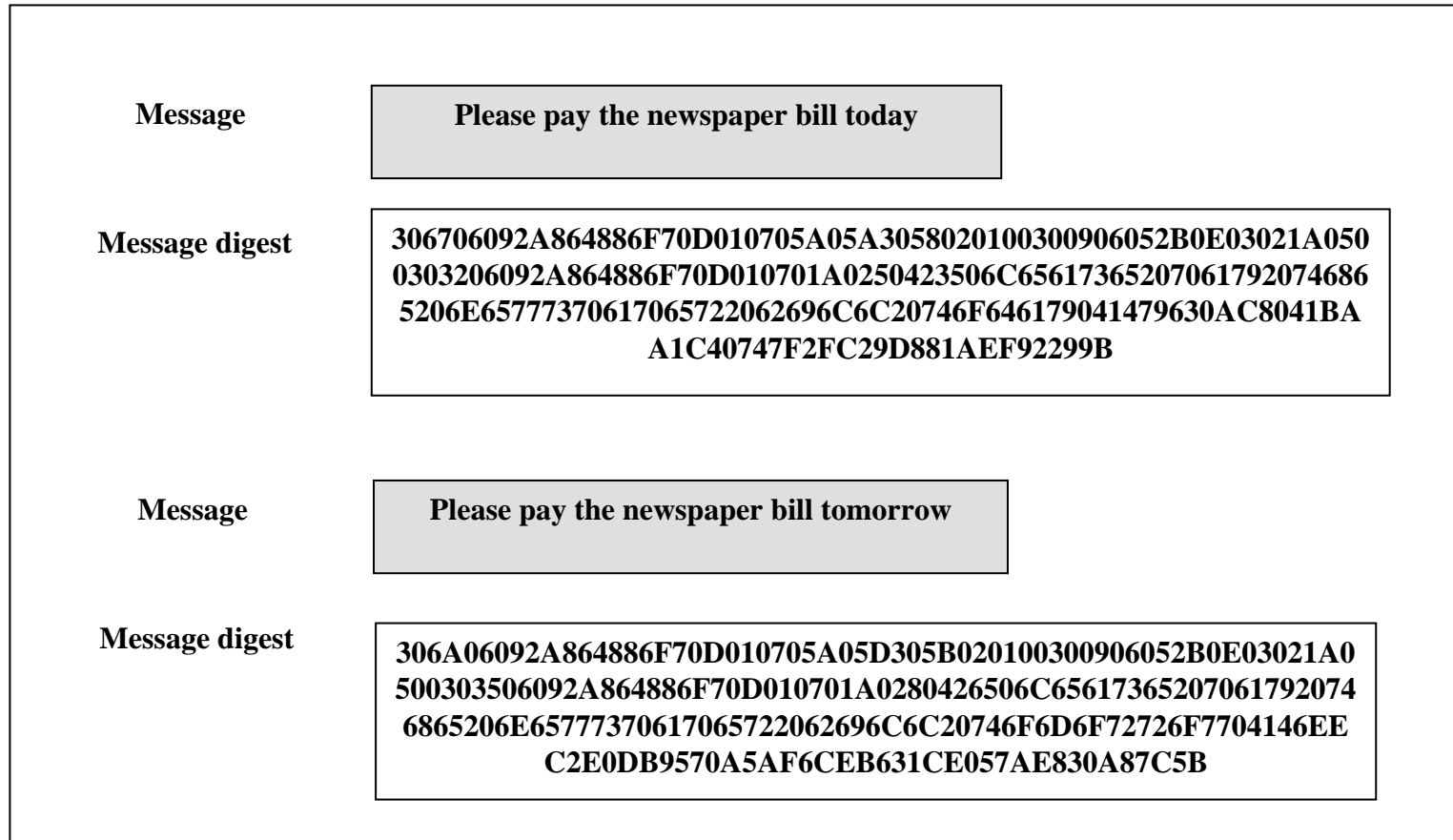


Fig 4.23

Message Digest Algorithms

- Basic principle: Take the original message, and reduce it to a smaller *fingerprint*
- Examples: MD5, SHA-1
- SHA-1 is considered stronger

One MD5 Operation

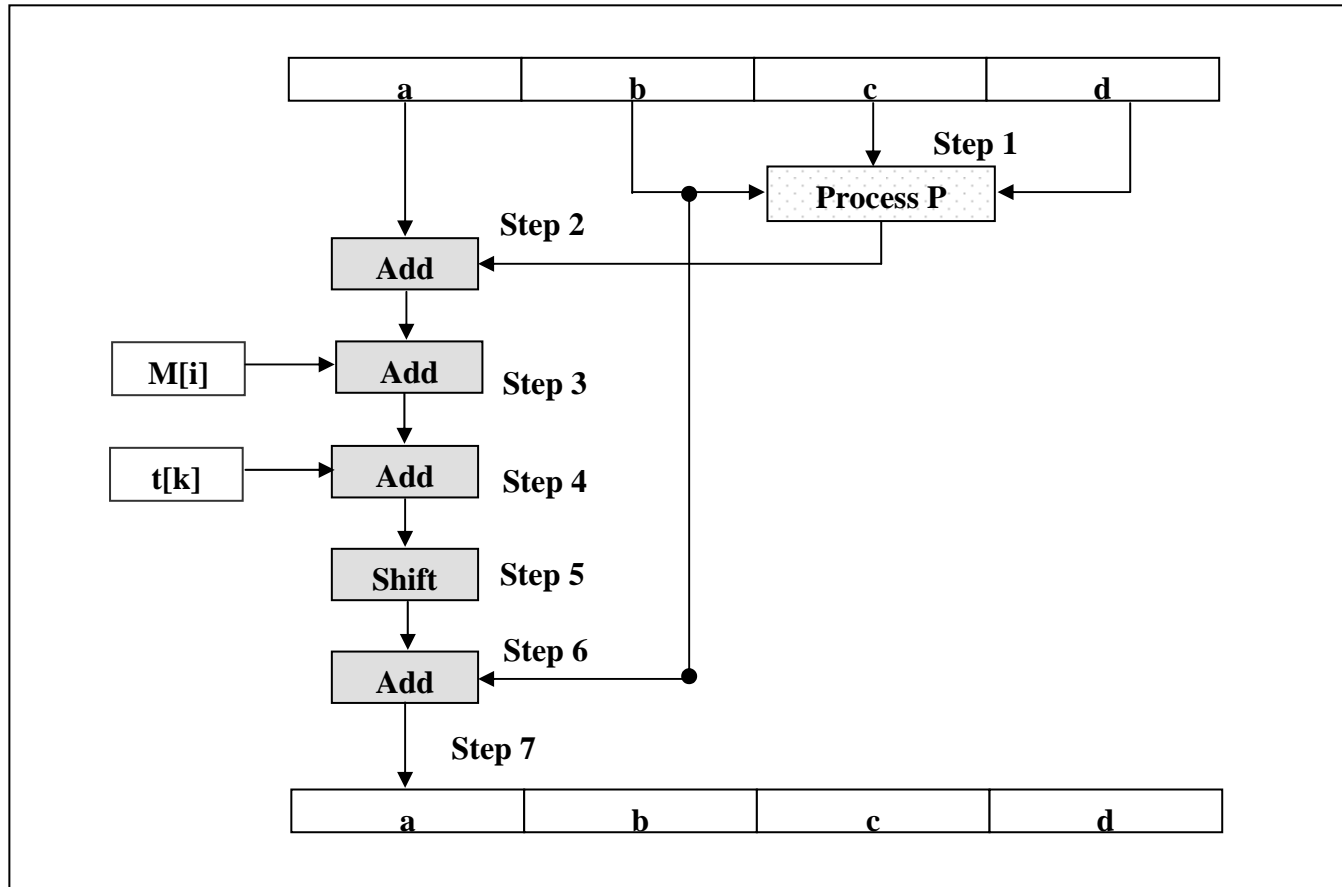


Fig 4.33

Single SHA-1 Iteration

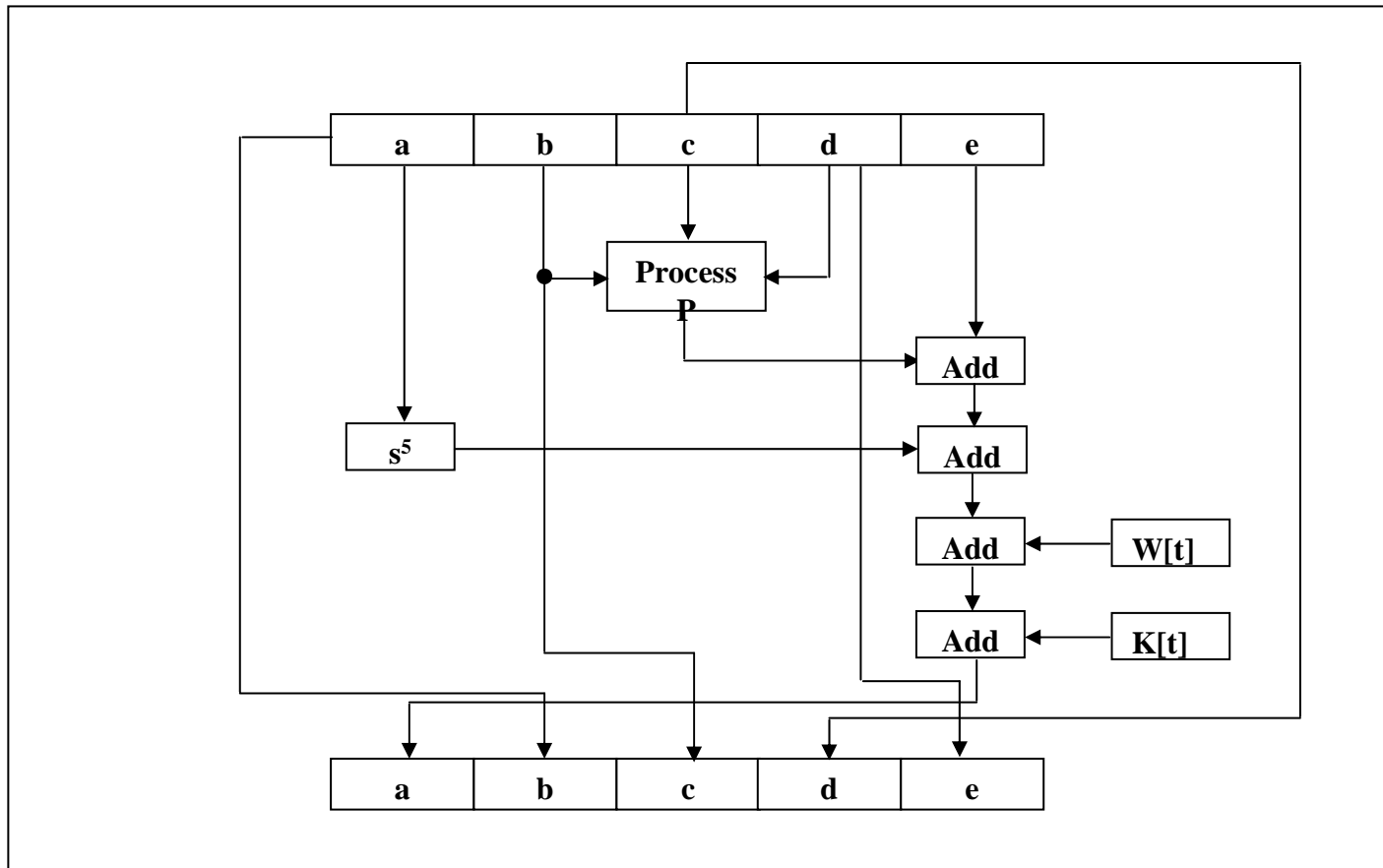


Fig 4.39

Comparison of MD5 and SHA-1

Point of discussion	MD5	SHA
Message digest length in bits	128	160
Attack to try and find the original message given a message digest	Requires 2^{128} operations to break in	Requires 2^{160} operations to break in, therefore more secure
Attack to try and find two messages producing the same message digest	Requires 2^{64} operations to break in	Requires 2^{80} operations to break in
Successful attacks so far	There have been reported attempts to some extent (as we discussed earlier)	No such claims so far
Speed	Faster (64 iterations, and 128-bit buffer)	Slower (80 iterations, and 160-bit buffer)
Software implementation	Simple, does not need any large programs or complex tables	Simple, does not need any large programs or complex tables

Fig 4.42

Message Authentication Code (MAC)

- Similar to message digest
- In addition, also involves encryption and decryption
- Sender and receiver must know a shared secret key

Message Authentication Code (MAC)

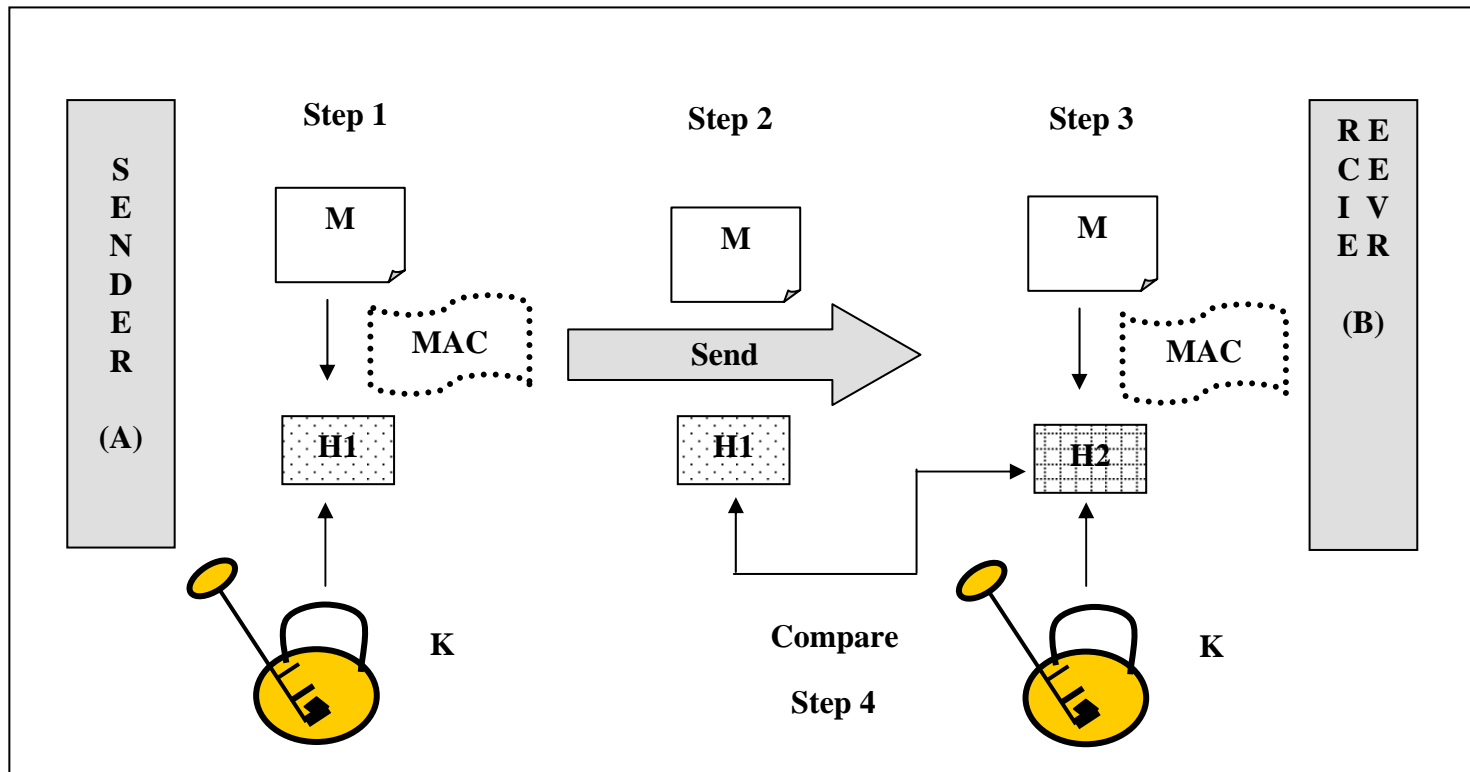


Fig 4.43

HMAC Concept

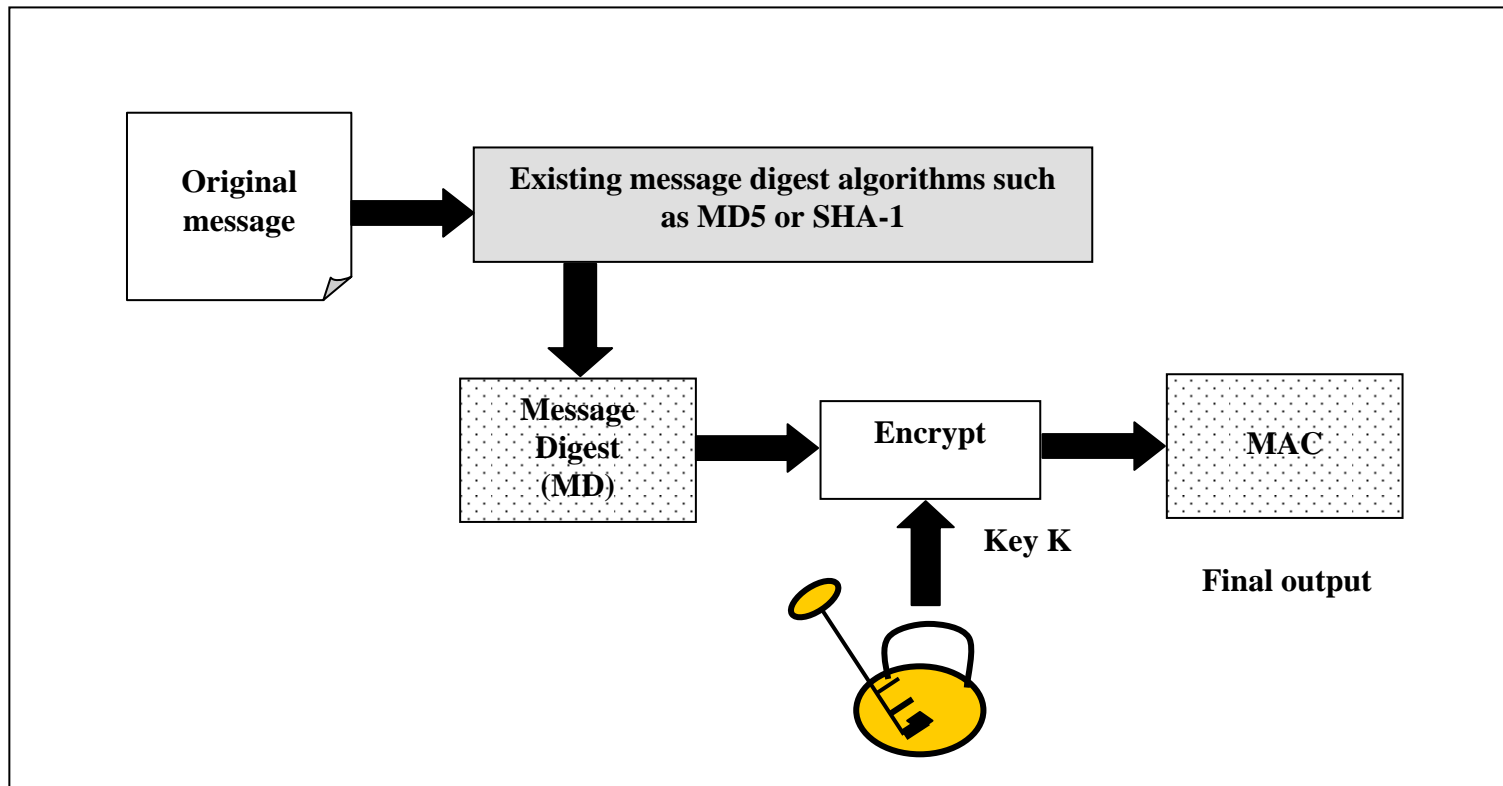


Fig 4.44

Complete HMAC Operation

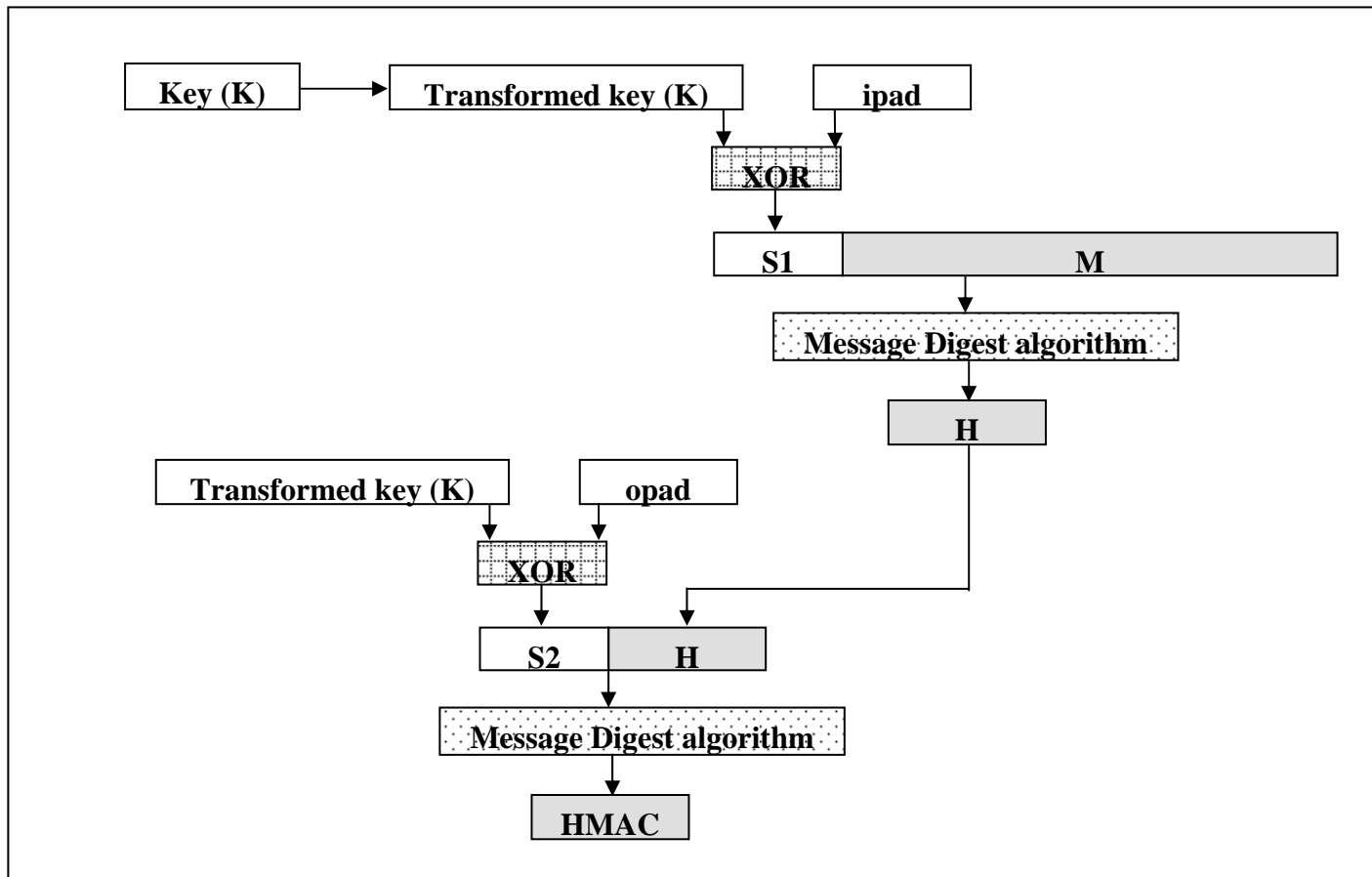


Fig 4.52